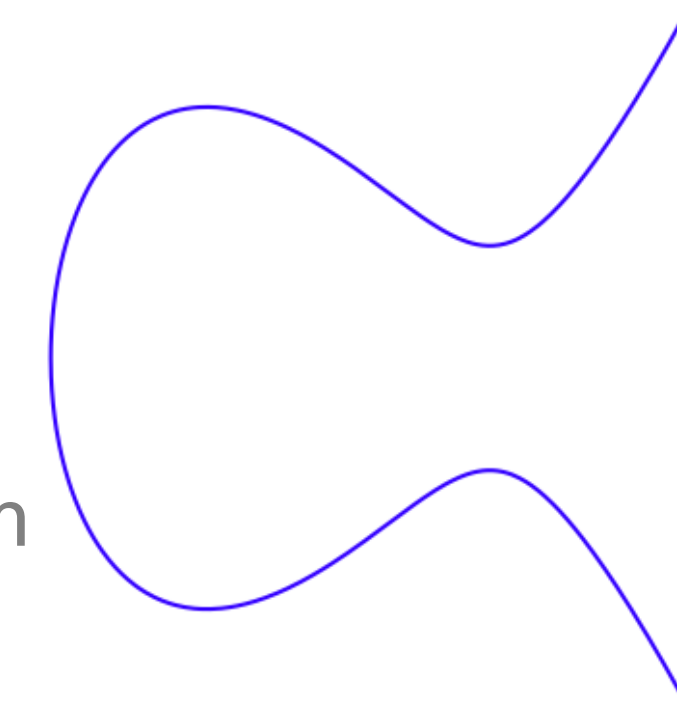


Towards a database of isogeny graphs

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The isogeny graphs $\Gamma_1(\ell; p)$

- An elliptic curve E over a finite field F_q has equation $E: y^2 = x^3 + Ax + B$: The j -invariant

$$j(E) = 1728 \frac{4A^3}{4A^3 + 27B^2}$$

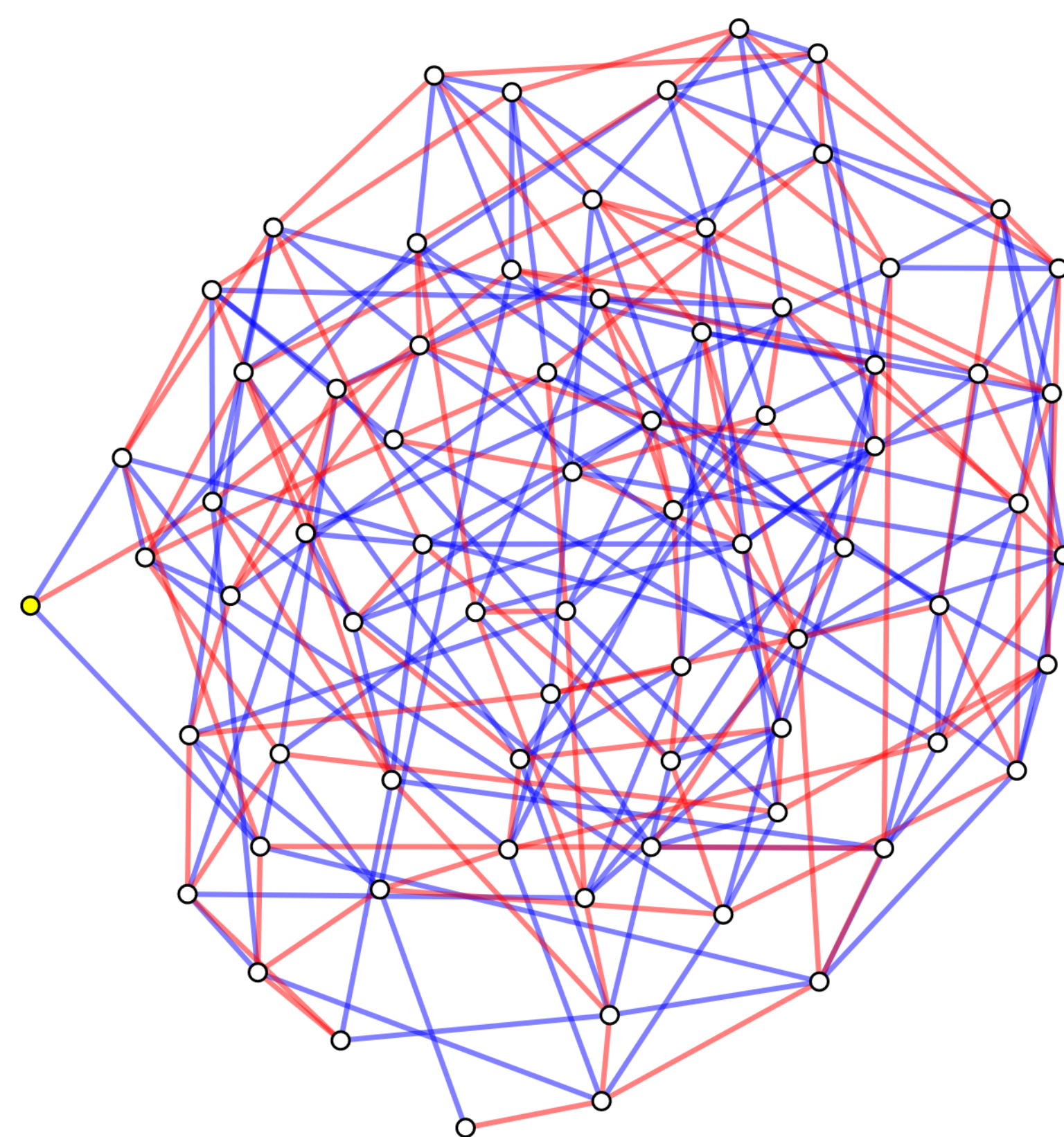
gives the isomorphism class of the curve. The curve is supersingular if its endomorphism ring is an order in a quaternion algebra.

- An isogeny of degree ℓ is a finite morphism of elliptic curves

$$\ell: E \rightarrow E^\ell$$

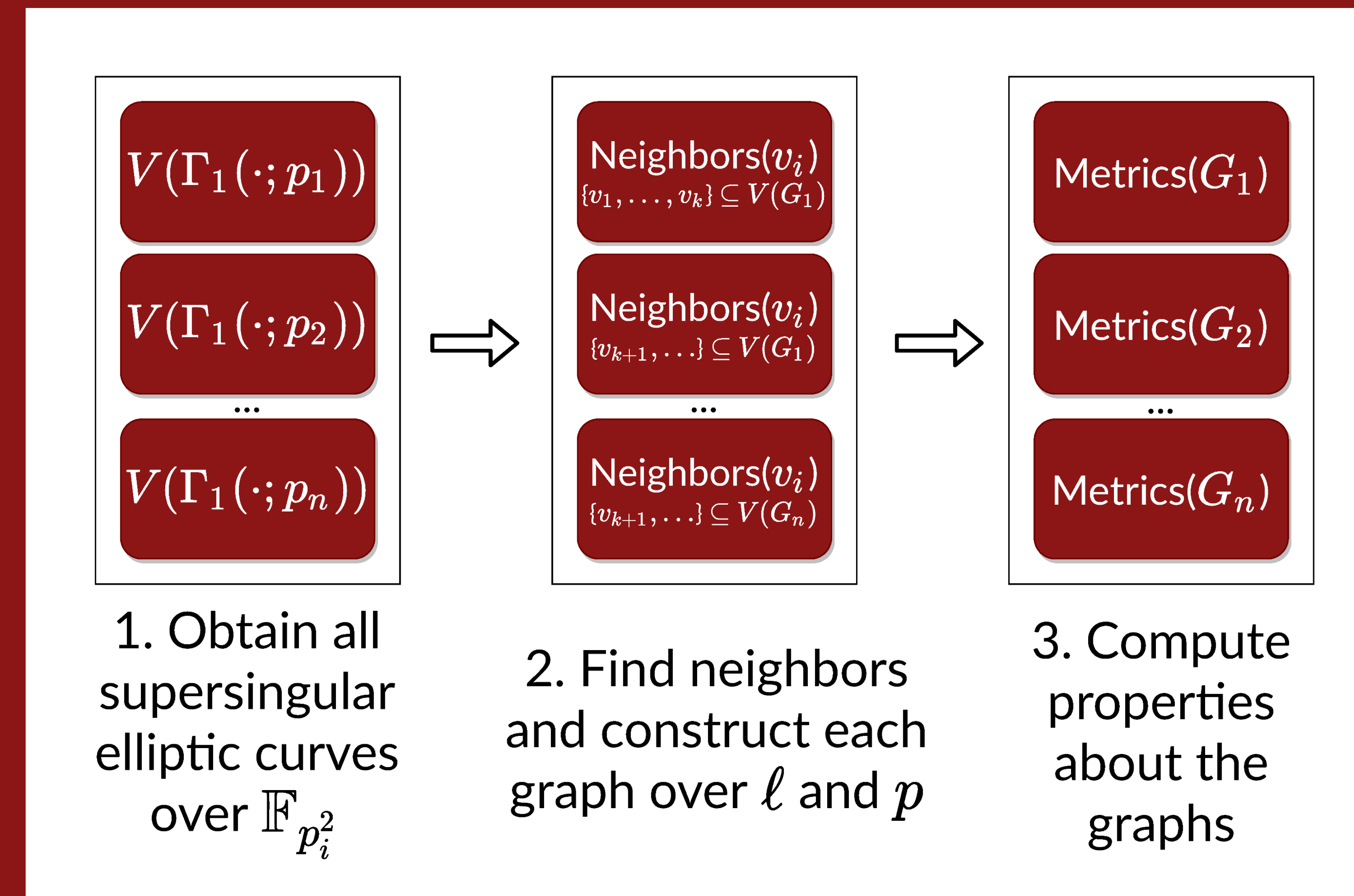
such that $j(\ker \ell) = j$.

- The supersingular isogeny graph $\Gamma_1(\ell; p)$ is the graph of supersingular j -invariants over F_p and degree- ℓ isogenies between them. It is an $(\ell + 1)$ -regular Ramanujan graph.



$\Gamma_1(2; 863)$ and $\Gamma_1(3; 863)$ superimposed.

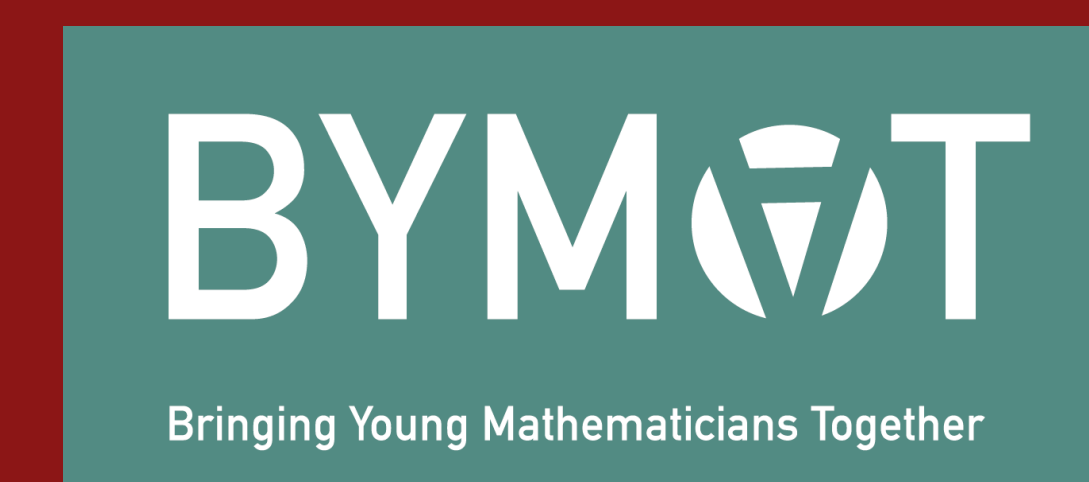
We are building an open library of supersingular isogeny graphs of elliptic curves using serverless computing



Each step is computed in parallel



Visit the website:
isogenies.enricflorit.com

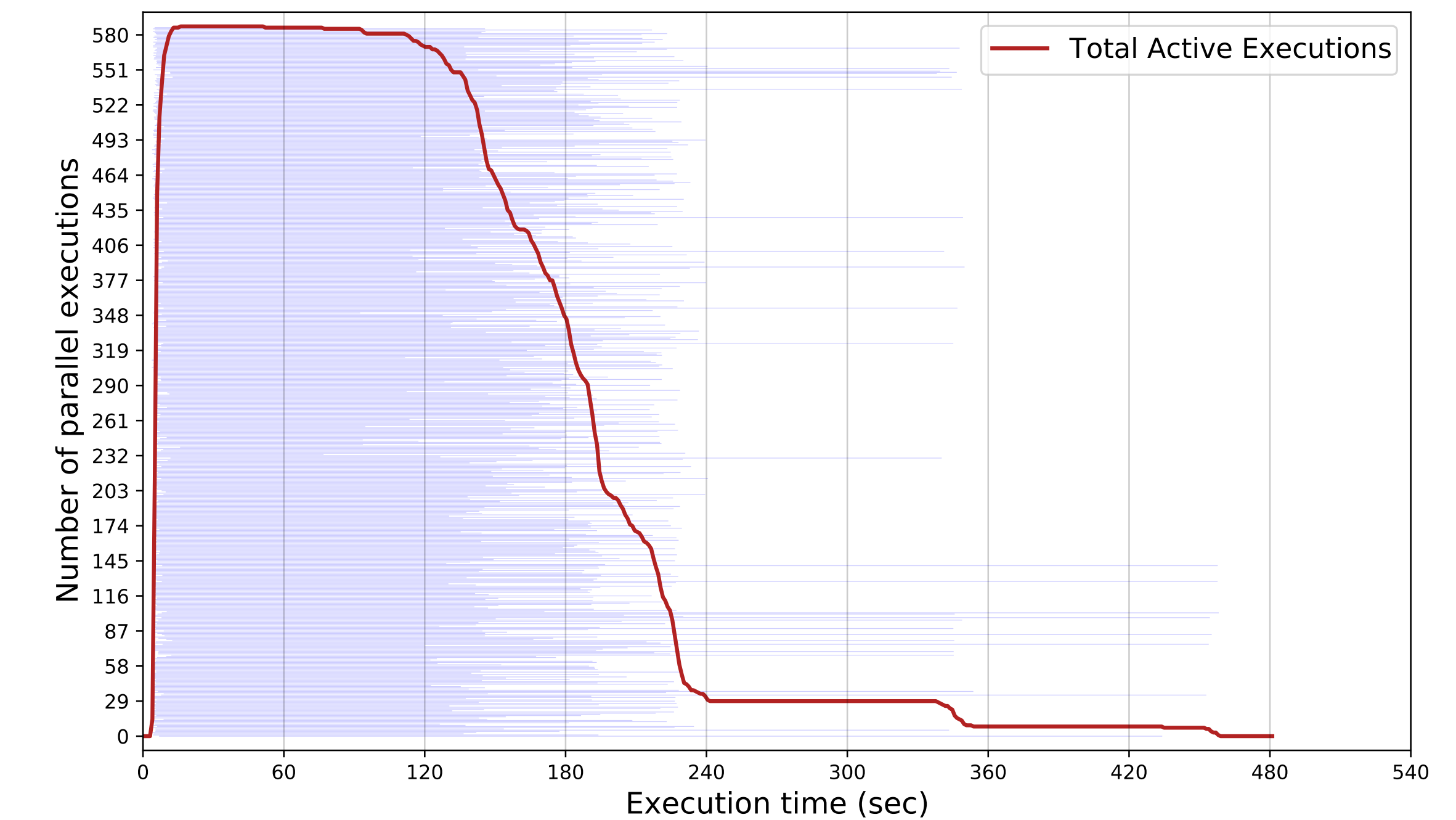


Computing isogenies

For each prime ℓ , the ℓ th modular polynomial $\Phi_\ell(X; Y) \in \mathbb{Z}[X; Y]$ satisfies $\Phi_\ell(j(E); j(E^\ell)) = 0$ ($\ell > 3$) there is a degree- ℓ isogeny from E to E^ℓ . The roots of $\Phi_\ell(j; Y)$ are the neighbors of j .

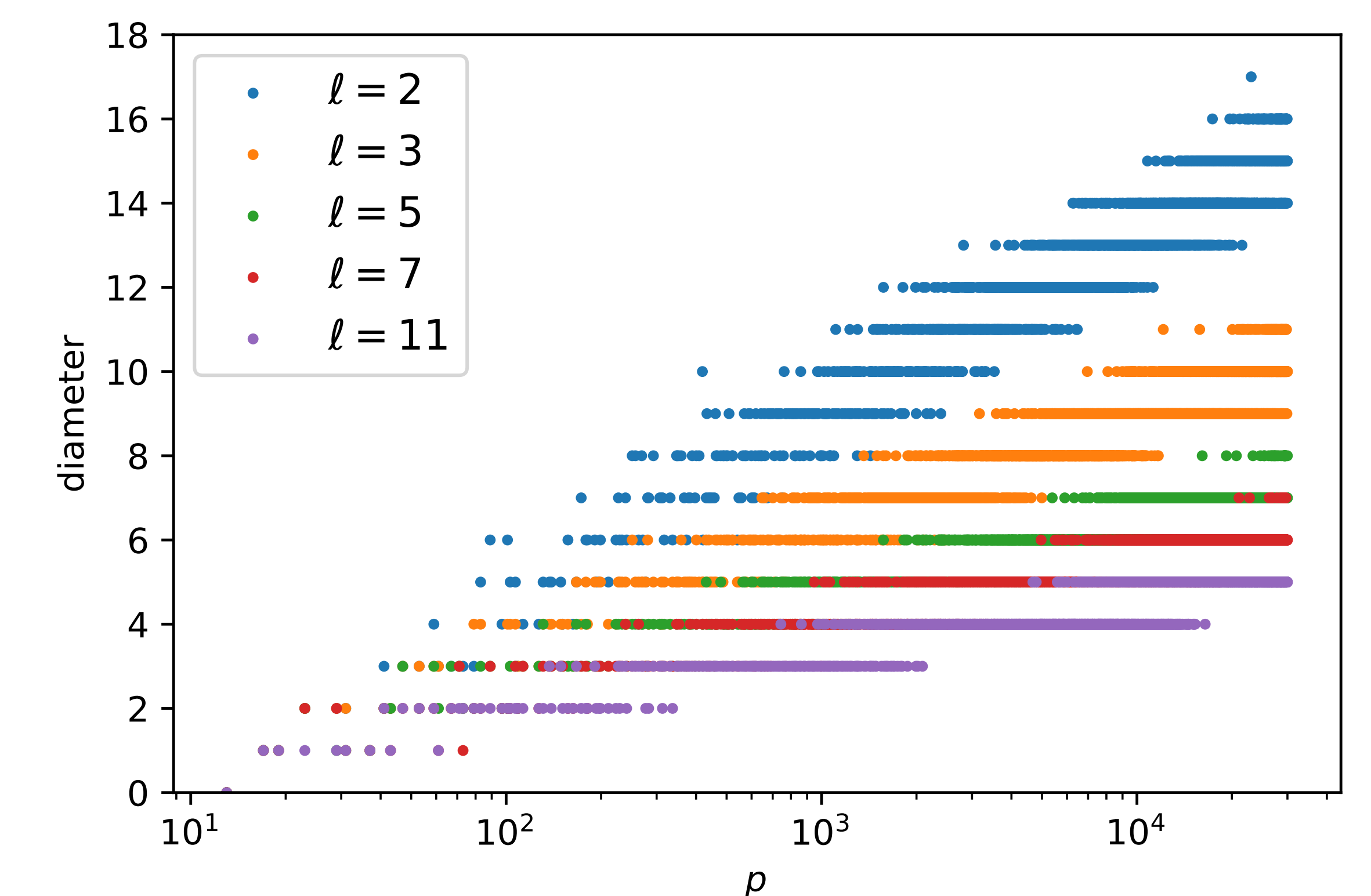
Serverless computing with Lithops

Using Lithops we can run hundreds of parallel distributed threads to compute graphs.



Histogram of the computation of $\Gamma_1(11; 4010173)$, overall time of 458 seconds.

Graph diameters



The diameter of $\Gamma_1(\ell; p)$ grows as $\log p$, as predicted by the Ramanujan property.

Main references

- [1] S. Arpin et al., "Adventures in Supersingularland", *arXiv preprint arXiv:1909.07779* 2019.
- [2] J. Sampe et al., "Towards Multicloud Access Transparency in Serverless Computing", *IEEE Software* 2020, DOI 10.1109/MS.2020.3029994.